# Improved Reduction Rules, Implemented in Peter's Engine 

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The rules are presented in Figures 1 and 2 and they are relative to a given initial marking $M_{0}$ and a cardinality query $\varphi$, where $\operatorname{places}(\varphi)$ is the set of all places that occur in the query $\varphi$.

Theorem 1. Let $\left(N, M_{0}\right)$ be a marked Petri net and let $\varphi$ be a cardinality query. Let $N^{\prime}$ be the net $N$ after the application of some reduction rules from Figures 1 and 2. Then $\left(N, M_{0}\right) \models E F \varphi$ if and only if $\left(N^{\prime}, M_{0}\right) \models E F \varphi$.

Theorem 2. Let $\left(N, M_{0}\right)$ be a marked Petri net. Let $N^{\prime}$ be the net $N$ after the application of some reduction rules from Figures 1 and 2 for a query $\varphi=2<1$. Then $\left(N, M_{0}\right)$ has a deadlock if and only if $\left(N^{\prime}, M_{0}\right)$ has a deadlock.

For the inhibitor-arc, we use $I(p, t) \in \mathbb{N} \cup\{0\}$. As a shorthand we write $I(p)=\{t \mid t \in T$ and $I(p, t) \neq 0\}$ (and $I(t)=\{p \mid p \in P$ and $I(p, t) \neq 0\})$ to denote the set of transitions (or places) which are connected via inhibitor-arcs.



Conditions on $p, t$ and $p^{\prime}$ :

$$
\begin{aligned}
& \text { A1: } p \subseteq p^{\bullet} \\
& \text { A2: } I(p)=\emptyset \\
& \text { A3: } p \notin \text { places }(\varphi) \\
& \text { A4: for all } t \in p^{\bullet} \text { it holds that } \\
& \\
& M_{0}(p)<F(p, t) .
\end{aligned}
$$

UE1: Remove $p$
UE2: Remove all $t \in p^{\bullet}$
(c) Rule E: Dead place removal

Fig. 2: Parallel rules for a cardinality formula $\varphi$ and initial marking $M_{0}$

